

Lecture 8

Gidon Rosalki

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1 Reminder

Last lecture we discussed HMMs (Hidden Markov Models):

$$\begin{aligned}\mathbb{P}[X_1 \dots X_n, H_1 \dots H_n] &= \mathbb{P}_0[H_0] \cdot \prod_i \mathbb{P}[H_{i+1}|H_i] \cdot \prod_i \mathbb{P}[X_i|H_i] \\ F_i[s] &= \mathbb{P}[H_i = s, X_1 \dots X_i] \\ B_i[s] &= \mathbb{P}[X_{i+1} \dots X_n | H_i = s]\end{aligned}$$

We also discussed likelihood, with the probability of a chain being

$$\begin{aligned}\mathbb{P}[X_1 \dots X_n] &= \sum_s F_n[s] \\ &= \sum_s B_1[s] \mathbb{P}_1[s] \\ \mathbb{P}[H_i = s | X_1 \dots X_n] &\propto F_i[s] \cdot B_i[s] \\ \mathbb{P}[H_i = s | H_{i+1} = t | X_1 \dots X_n] &\propto \text{something, we will discuss soon}\end{aligned}$$

2 Maximum probability reconstruction

We want to find the most likely Markov Chain

$$\arg \max_{h_1 \dots h_n} \{\mathbb{P}[H_1 = h_1 \dots H_n = h_n, X_1 \dots X_n]\}$$

How do we solve this problem? Dynamic Programming! I'm shocked! We will construct the following expression, to help us reach a problem solvable with DP:

$$\begin{aligned}& \max_{h_1 \dots h_i} \{\mathbb{P}[H_1 = h_1 \dots H_i = h_i, X_1 \dots X_i]\} \\ & \max_{h_1 \dots h_{i-1}} \left\{ \max_{h_i} \{\mathbb{P}[H_1 = h_1 \dots H_{i-1} = h_{i-1}, X_1 \dots X_{i-1}] \cdot \tau[h_i, h_{i-1}] \cdot \pi[x_i, h_i]\} \right\} \\ & = \max_{h_1 \dots h_{i-1}} \left\{ \mathbb{P}\left[H_1 = h_1, \dots, H_{i-1} = h_{i-1}, X_1 \dots X_{i-1}\right] \left(\max_{h_i} \{\tau[h_i, h_{i-1}] \pi[x_i, h_i]\} \right) \right\}\end{aligned}$$

Since

$$\begin{aligned}\mathbb{P}[H_1 \dots H_i, X_1 \dots X_i] &= \mathbb{P}[H_1 \dots H_{i-1}, X_1 \dots X_{i-1}] \\ &\quad \cdot \mathbb{P}[H_i | H_1 \dots H_{i-1}, X_1 \dots X_{i-1}] \\ &\quad \cdot \mathbb{P}[X_i | H_i, H_1 \dots H_{i-1}, X_1 \dots X_{i-1}]\end{aligned}$$

The final solution for the optimum comes out to

$$V_i[s] = \max_t \{V_{i-1}[t] \cdot \tau[s, t] \pi[x_i, s]\}$$

So, the optimal solution, of length i , finishing at the point $H_i = s$, is the optimal solution until $H_{i-1} = t$, of length $i - 1$, with the optimum transition between s and t .

Why are we using V ? Due to the name of the algorithm being **Viterbi**, named after its creator Andrew Viterbi.

Consider if you have a sequence $X_1 \dots X_n$, upon which you run Viterbi, and get $\hat{h}_1 \dots \hat{h}_n$. So, if we consider what happens at i :

$$\mathbb{P}[H_i = \hat{h}_i | X_1 \dots X_n]$$

What can we say about this?

3 Learning HMMs

Let us consider the full data

$$\begin{bmatrix} h_1 & h_2 & \dots & h_{n_1} \\ x_1 & x_2 & \dots & x_{n_2} \end{bmatrix}$$

So, firstly remember that each state is i.i.d. Then the probability of the data:

$$\begin{aligned} \mathbb{P}[\text{Data}] &= \prod_m \mathbb{P}[h_1[m] \dots h_{n_m}[m], X_1[m] \dots X_{n_m}[m]] \\ &= \prod_m \left[p_0[h_1[m]] \prod_{i=1}^{n_m-1} \mathbb{P}[H_i = h_i[m] | H_{i-1} = h_{i-1}[m]] \right] \\ &= \prod_m \left[\mathbb{P}_0[h_1[m]] \prod_{i=1}^{h_m-1} \tau[h_i[m], h_{i-1}[m]] \cdot \prod_i \pi[X_i[m], h_i[m]] \right] \end{aligned}$$

Let us define a statistic that counts how many times we had s

$$N_{0,s} = \sum_m \mathbb{1}_{n_1[m]=s}$$

Where

$$\begin{aligned} N_{t,s} &= \sum_m \sum_i \mathbb{1}_{h_{i-1}[m]=t, h_i[m]=s} \\ N_{s,x} &= \sum_r \sum_i \mathbb{1}_{h_i[r]=s, X_i[r]=x} \end{aligned}$$

So:

$$\begin{aligned} \mathbb{P}[\text{Data}] &= \prod_m \mathbb{P}[h_1[m] \dots h_{n_m}[m], X_1[m] \dots X_{n_m}[m]] \\ &= \left[\prod_s \mathbb{P}_0[s]^{N_{0,s}} \right] \cdot \left[\prod_{t \in S} \prod_{s \in S} \tau[s, t]^{N_{t,s}} \right] \cdot \left[\prod_{s \in S} \prod_{x \in \mathcal{X}} \pi[x, s]^{N_{s,x}} \right] \end{aligned}$$

Where \mathcal{X} is all the observations (ie, all the options that can happen in π).

We can now also discuss for our MLE:

$$\begin{aligned} \hat{p}_0[s] &= \frac{N_{0,s}}{\sum_t N_{0,t}} \\ \hat{\tau}[s, t] &= \frac{N_{t,s}}{\sum_u N_{t,u}} \\ \hat{\pi}[s, x] &= \frac{N_{s,x}}{\sum_y N_{s,y}} \end{aligned}$$

Let us now consider the case where we have many sequences $[X_1[1], \dots, X_{n_1}[1]]$, and $[X_1[2], \dots, X_{n_2}[2]]$. Our hidden state is now hidden. So

$$\begin{aligned} \mathbb{P}[\text{Data}] &\stackrel{iid}{=} \prod_m \mathbb{P}[X_1[m] \dots X_{n_m}[m]] \\ &= \prod_m \sum_{n_i} \sum_{h_{n_m}} \mathbb{P}[\text{same contents}] \end{aligned}$$

So this is hard, because combining together sums and multiplications is difficult. We cannot simply extract the sums from each other. We also do not know *anything* about the hidden states, there could be many between every other state.